## Cryptographic Protocols

## Exercise 12

### 12.1 Hyper-Invertible Matrices

Recall the definition of hyper-invertible matrices from the lecture:
Definition 1. A $r \times c$-matrix $M$ over some field $\mathbb{F}$ is called hyper-invertible if every square sub-matrix $M_{R}^{C}$ of $M$ is invertible, where, for sets $R \subseteq\{1, \ldots, r\}$ and $C \subseteq\{1, \ldots, c\}$ with $|R|=|C|>0, M_{R}^{C}$ denotes the matrix consisting of rows $i \in R$ and colums $j \in C$ of $M$.
a) Determine whether the following matrices are hyper-invertible:

$$
\left.\begin{array}{ll}
A & =\left[\begin{array}{lllll}
5 & 4 & 3 & 2 & 1 \\
1 & 2 & 3 & 5 & 5
\end{array}\right] \text { over } \operatorname{GF}(7)
\end{array} \quad B=\left[\begin{array}{lll}
4 & 1 & 4 \\
6 & 4 & 1 \\
3 & 1 & 1
\end{array}\right] \text { over } \operatorname{GF}(7) ~ 子 \begin{array}{lll}
1 & 2 & 3 \\
3 & 4 & 7 \\
4 & 2 & 6 \\
2 & 2 & 4
\end{array}\right] \text { over } \operatorname{GF}(11) \quad D=\left[\begin{array}{ccccc}
5 & 1 & 10 & 6 & 1 \\
1 & 6 & 0 & 1 & 5 \\
5 & 9 & 1 & 4 & 4 \\
4 & 7 & 5 & 5 & 2
\end{array}\right] \text { over GF(11) }
$$

Next, we want to show that permuting hyper-invertible matrices and multiplying columns (or rows) by constants preserves hyper-invertibility. Let $M \in \mathbb{F}^{r \times c}$ be a hyper-invertible matrix over some field $\mathbb{F}$.
b) Let $M^{\prime}$ be the matrix obtained from $M$ by exchanging the $i$ th and $j$ th column of $M$. Show that $M^{\prime}$ is hyper-invertible.
c) Let $\bar{M}$ be the matrix obtained from $M$ by multiplying each entry of the $i$ th column of $M$ by some value $a \in \mathbb{F} \backslash\{0\}$.
Show that $\bar{M}$ is hyper-invertible.

### 12.2 Properties of Hyper-Invertible Matrices

In this task we prove one direction of the lemma from the lecture: for a matrix $M$, which induces a linear function $f$, we have that $M$ is hyper-invertible if and only if $f$ is hyper-invertible.

Recall the definition of hyper-invertible mappings:
Definition 2. Consider a function $f: \mathbb{F}^{c} \rightarrow \mathbb{F}^{r}$, as well as some arbitrary inputs $\left(x_{1}, \ldots, x_{c}\right)$ and the corresponding function values $\left(y_{1}, \ldots, y_{r}\right)=f\left(x_{1}, \ldots, x_{c}\right)$. The function $f$ is called hyper-invertible if for any sets $A \subseteq\{1, \ldots, c\}, B \subseteq\{1, \ldots, r\}$ with $|A|+|B|=c$, there exists a function $f^{\prime}: \mathbb{F}^{c} \rightarrow \mathbb{F}^{r}$ that maps the values $\left\{x_{i}\right\}_{i \in A},\left\{y_{i}\right\}_{i \in B}$ to the values $\left\{x_{i}\right\}_{i \in \bar{A}},\left\{y_{i}\right\}_{i \in \bar{B}}$.

Prove that any hyper-invertible matrix defines a hyper-invertible linear function.
Hint: Note that for $A, B$ as in Definition 2 we have $\vec{y}_{B}=M_{B} \vec{x}=M_{B}^{A} \vec{x}_{A}+M_{B}^{\bar{A}} \vec{x}_{\bar{A}}$.

### 12.3 Beaver's Multiplication Triples

In the lecture we saw a multiplication protocol based on precomputing random double sharings. An alternative multiplication protocol can be obtained by precomputing multiplication triples, which are sharings of values $(a, b, c)$, all shared by polynomials of degree $t$, where $a$ and $b$ are uniform random values and $c=a b$.
Let $(a, b, c)$ be a multiplication triple. Given a share of $[x]$ and a share of $[y]$, how can a party compute a share of $[x y]$ efficiently?

