ETH Zurich, Department of Computer Science SS 2021

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## **Cryptographic Protocols** Exercise 12

## 12.1 Hyper-Invertible Matrices

Recall the definition of hyper-invertible matrices from the lecture:

**Definition 1.** A  $r \times c$ -matrix M over some field  $\mathbb{F}$  is called *hyper-invertible* if every square sub-matrix  $M_R^C$  of M is invertible, where, for sets  $R \subseteq \{1, \ldots, r\}$  and  $C \subseteq \{1, \ldots, c\}$ with |R| = |C| > 0,  $M_R^C$  denotes the matrix consisting of rows  $i \in R$  and columns  $j \in C$ of M.

a) Determine whether the following matrices are hyper-invertible:

$$A = \begin{bmatrix} 5 & 4 & 3 & 2 & 1 \\ 1 & 2 & 3 & 5 & 5 \end{bmatrix} \text{ over } \operatorname{GF}(7) \quad B = \begin{bmatrix} 4 & 1 & 4 \\ 6 & 4 & 1 \\ 3 & 1 & 1 \end{bmatrix} \text{ over } \operatorname{GF}(7)$$
$$C = \begin{bmatrix} 1 & 2 & 3 \\ 3 & 4 & 7 \\ 4 & 2 & 6 \\ 2 & 2 & 4 \end{bmatrix} \text{ over } \operatorname{GF}(11) \qquad D = \begin{bmatrix} 5 & 1 & 10 & 6 & 1 \\ 1 & 6 & 0 & 1 & 5 \\ 5 & 9 & 1 & 4 & 4 \\ 4 & 7 & 5 & 5 & 2 \end{bmatrix} \text{ over } \operatorname{GF}(11)$$

Next, we want to show that permuting hyper-invertible matrices and multiplying columns (or rows) by constants preserves hyper-invertibility. Let  $M \in \mathbb{F}^{r \times c}$  be a hyper-invertible matrix over some field  $\mathbb{F}$ .

- b) Let M' be the matrix obtained from M by exchanging the *i*th and *j*th column of M. Show that M' is hyper-invertible.
- c) Let  $\overline{M}$  be the matrix obtained from M by multiplying each entry of the *i*th column of M by some value  $a \in \mathbb{F} \setminus \{0\}$ .

Show that  $\overline{M}$  is hyper-invertible.

## 12.2 Properties of Hyper-Invertible Matrices

In this task we prove one direction of the lemma from the lecture: for a matrix M, which induces a linear function f, we have that M is hyper-invertible if and only if f is hyper-invertible.

Recall the definition of hyper-invertible mappings:

**Definition 2.** Consider a function  $f : \mathbb{F}^c \to \mathbb{F}^r$ , as well as some arbitrary inputs  $(x_1, \ldots, x_c)$  and the corresponding function values  $(y_1, \ldots, y_r) = f(x_1, \ldots, x_c)$ . The function f is called *hyper-invertible* if for any sets  $A \subseteq \{1, \ldots, c\}, B \subseteq \{1, \ldots, r\}$  with |A| + |B| = c, there exists a function  $f' : \mathbb{F}^c \to \mathbb{F}^r$  that maps the values  $\{x_i\}_{i \in A}, \{y_i\}_{i \in \overline{B}}$ .

Prove that any hyper-invertible matrix defines a hyper-invertible linear function. HINT: Note that for A, B as in Definition 2 we have  $\vec{y}_B = M_B \vec{x} = M_B^A \vec{x}_A + M_B^{\overline{A}} \vec{x}_{\overline{A}}$ .

## 12.3 Beaver's Multiplication Triples

In the lecture we saw a multiplication protocol based on precomputing random double sharings. An alternative multiplication protocol can be obtained by precomputing *multiplication triples*, which are sharings of values (a, b, c), all shared by polynomials of degree t, where a and b are uniform random values and c = ab.

Let (a, b, c) be a multiplication triple. Given a share of [x] and a share of [y], how can a party compute a share of [xy] efficiently?