ETH Zurich, Department of Computer Science SS 2021

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## Cryptographic Protocols Exercise 3

## 3.1 Geometric Zero-Knowledge

In this exercise we consider geometric constructions using a ruler (without markings) and a compass (Lineal and Zirkel). The operations we consider are those that we know from high school, namely to draw a line through two points, to draw a circle with center at one point that goes through another point, to obtain the intersection between two lines/two circles/a line and a circle, and to copy circles.<sup>12</sup>

a) An *angle* is a geometric object consisting of two rays (half-lines) with a common end point. Show how one can add and subtract two angles, i.e., given angles  $\alpha$  and  $\beta$ , construct  $\alpha + \beta$  and  $\alpha - \beta$  using the above operations.

A well-known result from abstract algebra states that the trisection of an arbitrary angle cannot be drawn in the above sense.

- b) Peggy claims that she knows<sup>3</sup> the trisection  $\alpha$  of a publicly known angle  $\beta = 3\alpha$ . Construct an interactive protocol that allows her to prove this claim. You may assume that Peggy can generate a random point on a circle and that Vic can flip a fair coin.
- c) Prove that your protocol is complete and argue (informally) why it is a proof of knowledge.
- d) Prove that your protocol is zero-knowledge.

## 3.2 Honest-Verifier Zero-Knowledge and c-Simulatability

In this exercise, we investigate the relation between HVZK protocols and c-simulatable protocols. It is easy to see that any *c*-simulatable protocol is also HVZK, while the converse is generally not true.

However, one can prove that any 3-round HVZK protocol with challenge chosen uniformly from a challenge space C, can be transformed into a *c*-simulatable protocol. More concretely, given an HVZK protocol (P, V) for relation R (P proves knowledge of a witness w for an instance x such that  $(x, w) \in R$ ), find a *c*-simulatable protocol (P', V') for R.

## 3.3 An Interactive Proof

Let G be a cyclic group of prime order (i.e., |G| is prime) and let  $g, h \in G$  be publicly known.

- a) Construct a *zero-knowledge* interactive proof protocol that, given some  $z \in G$ , allows Peggy to convince Vic that she knows some pair (x, y) such that  $z = g^x h^y$ . Prove that your protocol is complete and argue (informally) why it is a proof of knowledge.
- b) Prove that your protocol is zero-knowledge.

<sup>&</sup>lt;sup>1</sup>This last operation can actually be performed with the other three.

<sup>&</sup>lt;sup>2</sup>If you desire, you may play with the applet on www.geogebra.org.

<sup>&</sup>lt;sup>3</sup>i.e., holds a copy of the geometric object  $\alpha$ .